

NoSQL data stores and SOS:

Uniform Access to Non-Relational Database Systems

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Outline

Context

- Relational DBMS
- NoSQL Data Stores
- NoSQL Timeline

NoSQL Data Stores

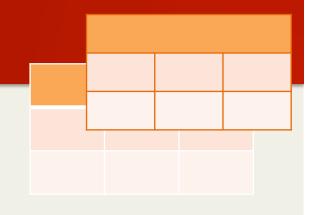
- Extensible Record Stores
- Document Stores
- Key-Value Stores
- Rise of NoSQL: a survey
- "NoSQL is about choice"
- Heterogeneity

• SOS - Save Our Systems

- Goal and requirements
- Data Model
- Interface
- Architecture
- Translation techniques

Future Work

Relational DBMS



Relational Databases

- Provide efficient support for applications that:
 - Require perstistence, consistence, avaliability, usability, ...
 - Simple data structure (numeric, string, ...)
 - Complex queries expressed by declarative languages

Relational DBMS

(Stonebraker & Cattel, CACM June 2011)

- "General-purpose traditional row stores"
 - Disk-oriented storage
 - Tables stored row-by-row on disk (hence, a row store)
 - B-trees as the indexing mechanism
 - Dynamic locking as the concurrency control mechanism
 - A write-ahead log (WAL) for crash recovery
 - SQL as the access language
 - A "row-oriented" query optimizer and executor

Relational DBMS

- Relational databases proposed as universal solution
 - It isn't completely true
 - OLAP, OLTP, XML, Stream processing, ...

One size does fit all?

(Michael Stonebraker, Ugur Çetintemel: "One Size Fits All": An Idea Whose Time Has Come and Gone. ICDE 2005: 2-11)

Why NoSQL?



- Many users
- Many concurrent interactions (reads and writes)
- Lots of data.

Need for:

- Scalable, distributed storage systems
- Flexible "web-proof" data models
- Easy interaction with programming languages



NoSQL Data Stores

Characteristics

- High scalability
- Data replication according to a distributed architecture
- Flexible data structure
- New indexing patterns

Missing features

- Simple interface
- New approach towards consistency

Consistency

- CAP Theorem: A distributed system cannot satisfy all the following properties:
 - Consistency
 - Availability
 - Partition-tolerance
- New Consistency Approaches
 - Strong consistency
 - Weak consistency
 - Eventual consistency

NoSQL Timeline

- **2006** Google BigTable
- **2007** *Amazon Dynamo*
- 2007 HBase
- 2008 Cassandra
- 2009 Voldemort, Redis, Riak, MongoDB

• • •

• **2011** – Oracle NoSQL

Extensible Record Stores

- Extensible Record Stores (also called column-family stores)
 - Relaxation of the Relational Model
 - Store tables of extensible records
 - Partitioned across multiple nodes
- **2006** Google BigTable
- **2007** *Amazon Dynamo*
- **2007** HBase

HBase

• HBase:

- Database made of Tables, containing Rows identified by a unique Id.
- Columns (named qualifiers) within a table are grouped into Column Families (CFs).
- Tables and CFs are static, defined in advance.
- Qualifiers are dynamic, can be added/removed at runtime.
- Within the same CF, different rows can have different qualifiers.

| Row id | Column families | | |
|----------|--|-----------------------|-------------------|
| | ColumnFamily1 | ColumnFamily2 | ColumnFamilyN |
| row_id_1 | qualifier1 = "value1" qualifier2 = "value2" | qualifier1 = "value1" | |
| row_id_2 | | | |
| | | | |
| row_id_m | | | |

HBase

| Row id | Column families | | |
|----------|--|-----------------------|-------------------|
| | ColumnFamily1 | ColumnFamily2 | ColumnFamilyN |
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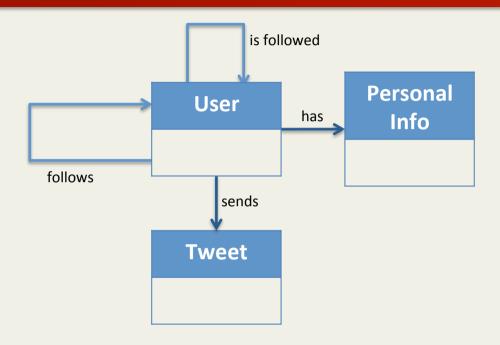
HBase:

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Unique data type:

- Byte-array
- Unstructured data stored in a semi-structured environment
 - Some assumptions are needed

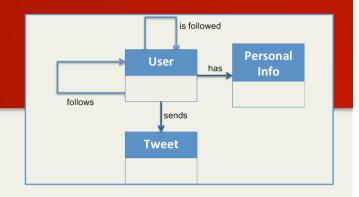
Example



Twitter Example

- Users, Tweets, Personal Info
- Users follows other users and are followed themself
- Users have personal info
- Users send tweets

Example



| Row id | Column families | | |
|--------|---|---|-------|
| | User Data | Personal Info | Tweet |
| 1001 | Username = "Alice" Password = TheAlicePassword | FirstName = "Alice" LastName = "Smith" | |
| 1002 | Username = "Bob" Password = TheBobPassword | | |
| | | | |

HBase

Operations

- Work on sigle rows or on lists of rows
- Provide direct access to rows given the row key
 - get(key)
 - put(key)
 - delete(key), deleteColumn(key, column), deleteFamily(...)
 - add(key, columnFamily), add(key, columnFamily, qualifier, value), ...
 - scan(table)
- Rows selection on the basis of filters defined on column families or qualifiers

— ...

HBase

Other characteristics

- Map-Reduce support (Hadoop)
- Strong consistency
- Max 10 column families
- Using filters deteriorates performances
- Bloom filters and column family compression for more efficient indexes

DynamoDB



DynamoDB

- Database made of Tables, containing Items identified by a unique key
- Items group a set of Attributes
- Attributes are characterized by a Name and a Value
- Every attribute can have multiple values
- Different Items belonging to the same table can have sets of disjoints attributes

Table

| Key | Other attributes |
|-----|---|
| key | (Name ₁ , value),, (Name ₁ , value) (Name ₃ , value) |
| key | (Name ₂ , value) |
| | |

DynamoDB

| Other attributes |
|---|
| (Name ₁ , value),, (Name ₁ , value) (Name ₃ , value) |
| (Name ₂ , value) |
| |

DynamoDB

- Database made of Tables, containing Items identified by a unique key
- Items group a set of Attributes
- Attributes are characterized by a Name and a Value
- Every attribute can have multiple values
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Provided on the cloud

Data types

- Scalar data types: number, string, binary
- Set data types: number set, string set, binary set

Example

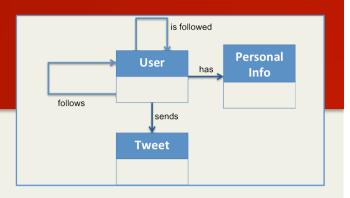


Table: Users

| Key | Other attributes |
|------|---|
| 1001 | Username = "Alice", Password = TheAlicePassword, FirstName = "Alice", |
| 1002 | Username = "Bob", Password = TheBobPassword, |
| ••• | |

DynamoDB

Operations

- work on sigle items
 - getItem(table, key)
 - putItem(table, key, av)
 - deleteItem(table, key)
- row selection on the basis of filters that use attribute names
- batch operations
 - batch putItem
 - batch deleteItem,
- ...

DynamoDB

Other characteristics

- High avaliability
- Keys are hashed: databases can be seen as a distributed hash table
- Node are located in specific regions (specified when the datastore is created)
- Data are replicated across nodes
- DynamoDB access cost policy
- Eventual consistency/strong consistenct

Document Stores

Document Stores

- Store collection of documents
- Documents are objects characterized by fields whose value can be a scalar, a list, a document itself.
- **2009** MongoDB

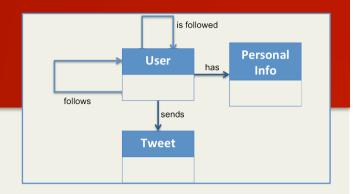
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MongoDB

- A Database is made of Collections
- A collection is a named group of **Documents**
- Documents are made of fields
- Fields value can be a scalar, a list, a document itself.

Users Collection

```
users: [
      _id: "1001",
       username: "bob1987",
       password: "TheBobPassword"
       personal: {
        firstName: "Bob",
        lastName: "Smith",
        ssn: "4hfe94"
      followers: [
           id: "2004",
           firstName: "Alice",
           lastName: "Smith",
           email: "alice@gmail.com"
           id: "1714",
```



Operations

- Operations defined on single fields:
 - insert(collection, doc)
 - find(selector, collection)
 - remove(selector, collection)
- Advanced operations
- The simplest selector is the empty document {} that matches all the documents of a collection.

Other characteristics

- Full index support
- Rich query API
- Sharding
- Strong consistency

Key-value datastores

Key-value data datastores

- Store values and an index for finding them based on programmerdefined key.
- A database is a collection of key-value pairs.
- 2009 Redis
- **2011** Oracle NoSQL
- •

Redis

- A Database is a schema-less collection of key-values pairs
- Key-value index

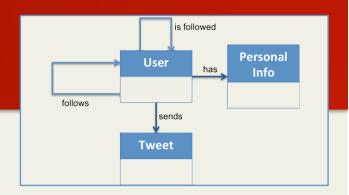
Data types

- Binary Strings: any type of binary data (byte array, number, plain string, ...).
- Integer counters
- Lists, Sets of binary strings
- Hashes

Complex data types cannot be nested

- It is not possible to define Sets of Lists, ...
- Unique key-space

```
users:1001:firstName = "Bob"
users:1001:lastName = "Smith"
users:1001 = {
 username = "bob1987"
  password = "TheBobPassword"
 friends:2004.email = "alice@gmail.com"
 friends:2004.firstName = "Alice"
```



Operations

- Simple operations:
 - set(key, value)
 - get(key)
 - delete(key)
- Advanced operations
 - Insert an element into a list or a set
 - Increment a counter
 - hgetall(key) that retrieves all the field-value pairs of a hash associated with the key

Other characteristics

- Efficient access (in memory)
- Map-Reduce support
- Strong consistency

NoSQL Data Stores

- Aspects to be considered
 - Number of accesses to retrieve an object
 - Resilience to unstructuredness
 - Partition-friendliness
 - How data are supposed to be queried
- Datastore best practices
- Performances influence data organization:
 - Denormalize data or not? Always? Never? When?
 - How we want to support consistency?

Rise of NoSQL: a survey

- 50% of IT managers/developers funded NoSQL projects in 2011
- 70% plan to fund NoSQL projects in 2012
- Enterprises in U.S.:
 - 56% already use some NoSQL database
 - 63% has plans to use in the next 2 years

"NoSQL is about choice"

- Jan Lenhardt (Couchbase Co-Founder)

Many data model families

- Key-value store
- Column-based store
- Document store
- Graph store

Many query models

- CRUD operations
- Map/Reduce queries
- Custom query languages
- Traversals

Many architectural choices

- Replicas (DHT?) vs sharding
- In-RAM vs traditional storage
- AP vs CP vs CA
- Strong vs eventual consistency
- **–** ...

"NoSQL is about choice" (2)

- Jan Lenhardt (Couchbase Co-Founder)

- Choose the right tool for your needs
 - One size does not fit all
- Do you need Map/Reduce?
 - Pick HBase or CouchDB
- Do you need great performances on simple operations?
 - Pick Redis
- Chances are you may need both

Heterogeneity still a problem

What if:

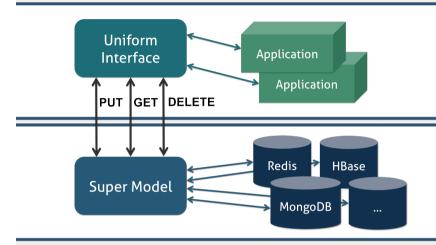
- I want to use many data stores at the same time
- I want to migrate my data
- I want to decouple my app from a specific technology

Reverse the canonical problem:

- One size (data store) does not fit all (apps)...
- ...but one size (your app) should fit all (the data stores)

- Goal: seamless access to different NoSQL data stores.
 - Define access
 - Define seamless
- Requirements:
 - Lightweight: small footprint on performances
 - Coherent: with main NoSQL themes and features
 - Hint: do not reimplement SQL
 - Seriously, someone has done it
 - Scalable: easily extendable to different technologies and data stores

SOS – Save our systems



Common Interface

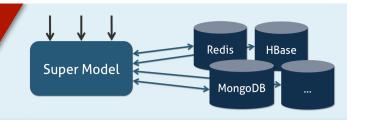
to access different NoSQL systems

Common Data Model

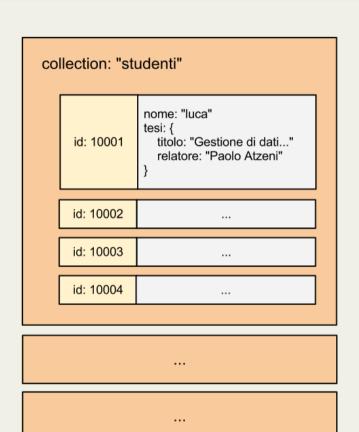
instances are mapped to the data stores of choice

- SOS is a Database Access Layer between the app and the data store
 - It collects data from the interface and seamlessly manages its translation and deployment to specific data stores
- Implementations provided for three data stores belonging to different families:
 - HBase (column-based store)
 - Redis (key-value store)
 - MongoDB (document store)

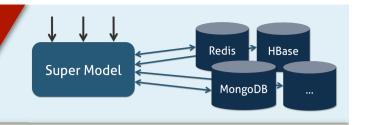
Common Data Model



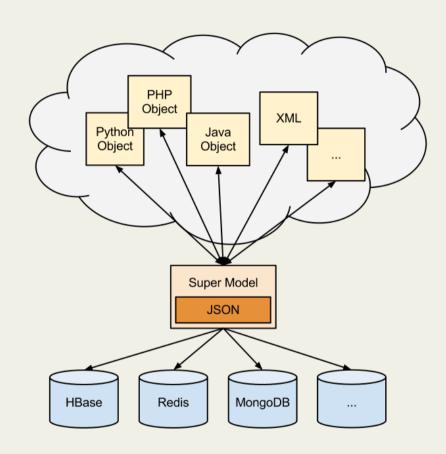
- SOS let users define collections of schemaless, tree-shaped objects
- Each collection is identified by a unique name
- Each object is identified by an ID, unique within the collection it belongs to



Common Data Model (2)

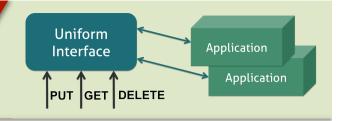


- Objects are materialized in JSON (JavaScript Object Notation)
 - Lightweight
 - Widely adopted
 - Platform-independent
- Custom translations are defined between JSON and each data store
- Translations are optimized to exploit efficiently the data store native structures



Common Interface

SOS – Save our Systems



Operations on single objects:

- put (collection-name, id, object) : void
- get (collection-name, id, type<T>) : <T>
- delete(collection-name, id) : void

Operations on single fields:

- put (collection-name, id, path, object) : void
- get (collection-name, id, path, type<T>) : <T>
- delete (collection-name, id, path) : void

Operations on collections:

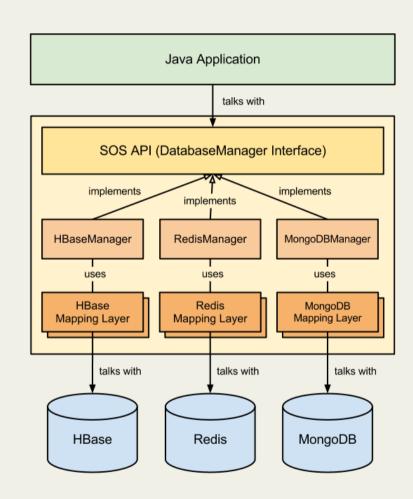
- get (collection-name, type<T>) : Collection<T>
- delete (collection-name) : void

Architecture

SOS – Save our Systems

• SOS is currently implemented as a **Java library**.

• It defines a **streamlined API** implemented by specific data store modules.



Usage example

SOS – Save our Systems

```
Student luca = new Student(...);
DatabaseHandler db = new HBaseHandler();
db.put("students", luca.getId(), luca);
```

public class Student {

private String id;

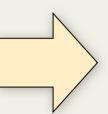
private String firstName; private String lastName;

private Set<Record> records;

Usage example

```
public class Student {
                                              private String id;
                                               private String firstName;
                                               private String lastName;
                                               private Set<Record> records;
Student luca = new Student(...);
DatabaseHandler db = new HBaseHandler();
db.put("students", luca.getId(), luca);
                            object
          collection
                                         object
                              id
            name
```

Translation example



```
id = "281283",
firstName = "Luca",
lastName = "Rossi",
records = [
       id = "10001",
       course = {
           id = "20001",
           name = "Databases 101"
       date = "2011/06/12",
       grade = "A"
   },
{
       id = "10002",
       course = {
           id = "20004",
           name = "Computer Vision",
       date = "2011/05/21",
       grade = "B"
                      JSON
```

Translation example

```
id = "281283",
firstName = "Luca",
lastName = "Rossi",
records = [
       id = "10001",
       course = {
           id = "20001",
           name = "Databases 101"
       date = "2011/06/12",
      grade = "A"
      id = "10002",
       course = {
           id = "20004",
           name = "Computer Vision",
       date = "2011/05/21",
      grade = "B"
                     JSON
```



| students (table) | |
|--|--|
| _top | records[] |
| <pre>id = "281283" firstName = "Luca" lastName = "Rossi"</pre> | [0].id = "10001" [0].course.id = "20001" [0].course.name = "Databases 101" [0].date = "2011/06/12" [0].grade = "A" |
| | [1].id = "10001" [1].course.id = "20004" [1].course.name = "Computer Vision" [1].date = "2011/05/21" [1].grade = "B" |
| HBase | |

Future work

SOS – Save our Systems

Architecture:

 Deploy SOS as a web application, exposing a REST interface that deals with JSON objects.

Translations:

 Enable custom translations, providing ways for the users to map structures into others, and so forth.

Interface:

- Provide support for the creation (and manteinance) of indexes
- Provide support for "update" operations that involve multiple nodes

Support for other DBMSes

Amazon DynamoDB and Oracle NoSQL are underway!

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- "NoSQL is about choice"
- Heterogeneity

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Questions?